



# MICROPROCESSOR SYSTEMS

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## Simple Microprocessor

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# Introduction

- After the invention of microprocessor, the size shape and price of digital computer have changed a lot
- The microprocessor is an IC (Integrated Circuit) that contains much of the processing capabilities of a computer
- It's a Large Scale Integration chip that is programmable
- A computer with a microprocessor as a CPU (Central Processing Unit) is called a microcomputer
- Microcomputer uses Stored Program concept
- It makes them reprogrammable

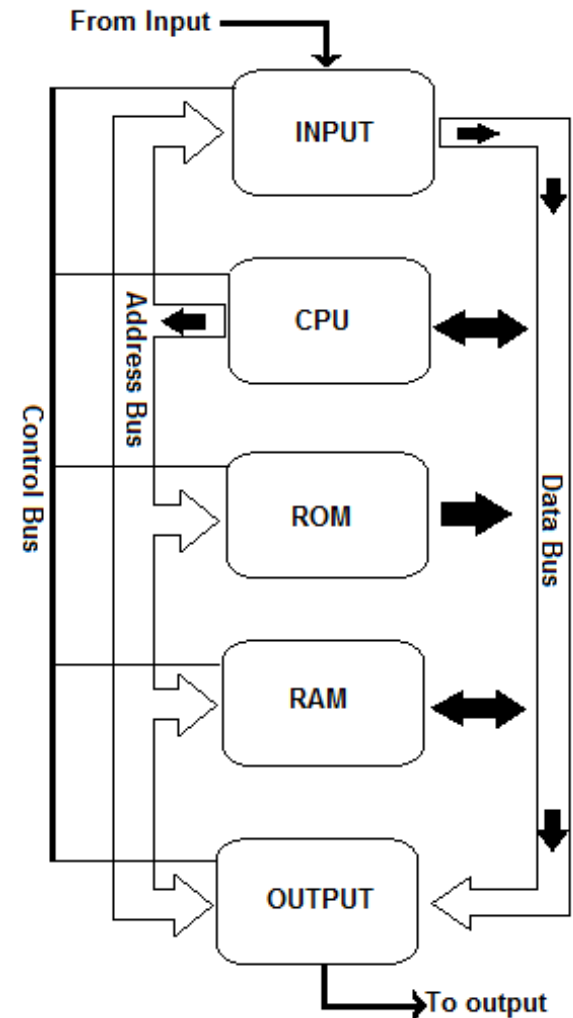


# A simple Microcomputer

- A digital microcomputer contains five functional components
- Input, Output, CPU, ROM and RAM (Hardware parts)
- ROM mostly stores program or list of instructions
- These programs manipulate data present in RAM
- Programs stores in ROM are called *firmware*
- The microprocessor controls all units with control lines
- Besides control lines:
  - Address lines (16) selects a memory or port location
  - Data lines (8) are two-way paths to transfer data to and from memory / port
- In this simplified organization details have been omitted

# A simple Microcomputer

- Two types of semiconductor memories (ROM and RAM)
- Input and output devices
- Bidirectional Data lines and Address lines





# Simplified CPU Organization and Instruction Execution

- Due to the wide variety of microprocessor vendors and microcomputer assembling companies, a generic microcomputer is suitable to study
- While learning about a new microprocessor a programmer must know about its:
  - Architecture
  - Instruction set
  - Minimal System
  - Control Signals
  - Pin functions

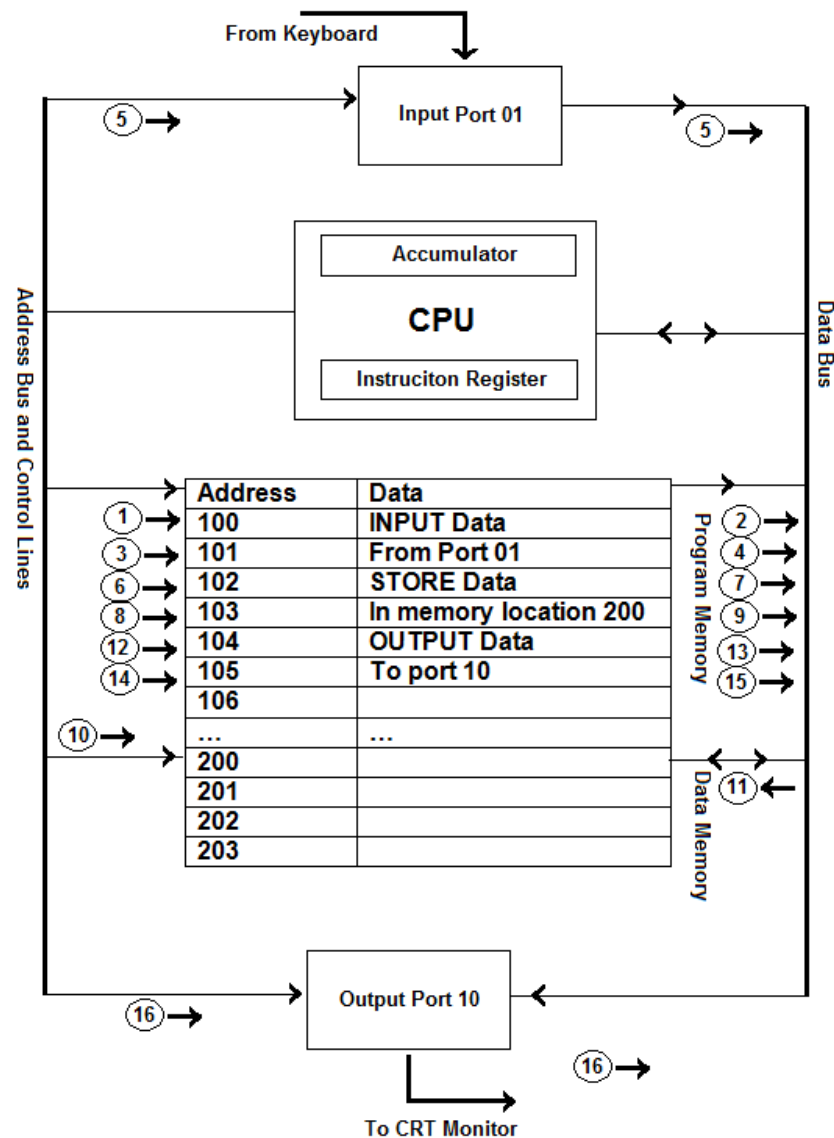


# Microcomputer Operations

- Consider the example of a simple microcomputer operation
  - Press a key on the keypad
  - Store the letter in memory
  - Display the character
- The instructions are stored in memory using instruction set
- Most of the instructions are stored in more than one memory locations
  - Like the first instruction will be composed of operation (input) and the port address where keypad is connected
  - Second instruction will be composed of operation and the memory address
  - Third instruction will be composed of operation and the display port number



# Microcomputer Operations Diagram





# Microcomputer Operations (Contd.)

- **Step 1:** CPU sends the address 100 on the address bus followed by enabling the *read* control line to memory
- **Step 2:** The memory puts the data byte on that address on the data bus. CPU reads it and put in a register (instruction register). CPU then decodes the instruction
- **Step 3:** CPU sends out the next address 101 to read the data for the first instruction followed by enabling the *read* control line to memory
- **Step 4:** The memory puts on the data bus the port number (01) from where to read the data. CPU can now interpret the whole instruction





# Microcomputer Operations (Contd.)

- **Step 5:** The CPU issues the address of port (01) on the address bus followed by read control line. The port then transfers the data on the data bus
- **Step 6:** The MPU addresses the location 102 on the address bus followed by read control signal
- **Step 7:** The operation code for STORE instruction is output by the memory and is brought inside the MPU. MPU decodes it
- **Step 8:** MPU sends next address 103 on the address bus to look for the operand followed by enabling read control line
- **Step 9:** Memory transfers the byte on the data bus. MPU can now decode the whole instruction



# Microcomputer Operations (Contd.)

- **Step 10:** (Execution) CPU now stores the data. It sends the address 200 on address bus and enable the write input to the memory
- **Step 11:** The MPU puts the data on the data bus.
- Step 12 to 16 are related to 3<sup>rd</sup> instruction



# Simplified Memory Architecture

- Random access and Sequential access
- The generic microprocessor has 16 address lines
- $2^{16}$  different combinations of 1's and 0's
- It is convenient to write them in hexadecimal
- So the first address is 0000h and last being FFFFh
- Generic microcomputer memory map
- First 256 bytes are ROM (1 page)
- In this organization memory is 256 X 4 bits
- There may be other organizations like 256 X 8bits, 256 X 1 bit
- The maximum memory supported by a microcomputer having 16 address lines is 64Kb

Address	Contents							
0000h	1	1	0	0	0	0	1	1
0001h	0	0	0	0	0	0	0	1
...								
00FEh	0	1	0	0	0	0	0	0
00FFh	1	0	0	0	1	1	1	0
0100h	1	0	0	0	0	0	0	1
...								
1FFFh	1	1	0	0	0	0	0	0
2000h	0	0	0	0	1	1	1	0
2001h	1	0	0	0	0	0	0	1
...								
20FEh	0	1	0	0	0	0	1	0
20FFh	1	0	0	0	0	0	0	1
2100h	0	1	0	0	0	0	1	1
...								
FFFFh	0	0	0	1	1	1	1	1



# Instruction Set

- The group of instructions a microprocessor can execute
- Generic microprocessor instruction set
- Instruction sets are not standardized
- Categories of instructions
  - Arithmetic Instructions (Addition, Subtraction, Shift, Increment, Decrement)
  - Logical Instructions (AND, OR, NOT, EX-OR etc.)
  - Data Transfer instructions (Memory to CPU, CPU to memory)
  - Branch Instructions (JUMP)
  - Others



# Arithmetic Instructions

- Following arithmetic operations
  - Addition between registers, registers and memory
  - Subtraction between registers, registers and memory
  - Increment register or memory contents
  - Decrement register or memory contents
  - Compare contents of two registers or memory – register
  - Negate the contents of register / memory contents



# Logical Instruction

- Includes:
  - Logical AND between registers or memory – register
  - Logical OR between registers or memory – register
  - Logical EX-OR between registers or memory – register
  - Logical NOT of register or memory word
  - Shift Left contents of register / memory
  - Shift Right contents of register / memory



# Data Transfer Instructions

- Includes:
  - Load a register with a value from memory (LOAD)
  - Store the register value to memory location (STORE)
  - Move data between registers (MOV)
  - Input from a port (INPUT)
  - Output to a port (OUTPUT)



# Branch Instructions

- Includes:
  - Unconditional Branch
  - Conditional Branch
    - Zero
    - Not Zero
    - Equal
    - Not Equal
    - Positive
    - Negative





# Other Instructions

- Includes:
  - Push and Pop
  - NOP
  - HLT



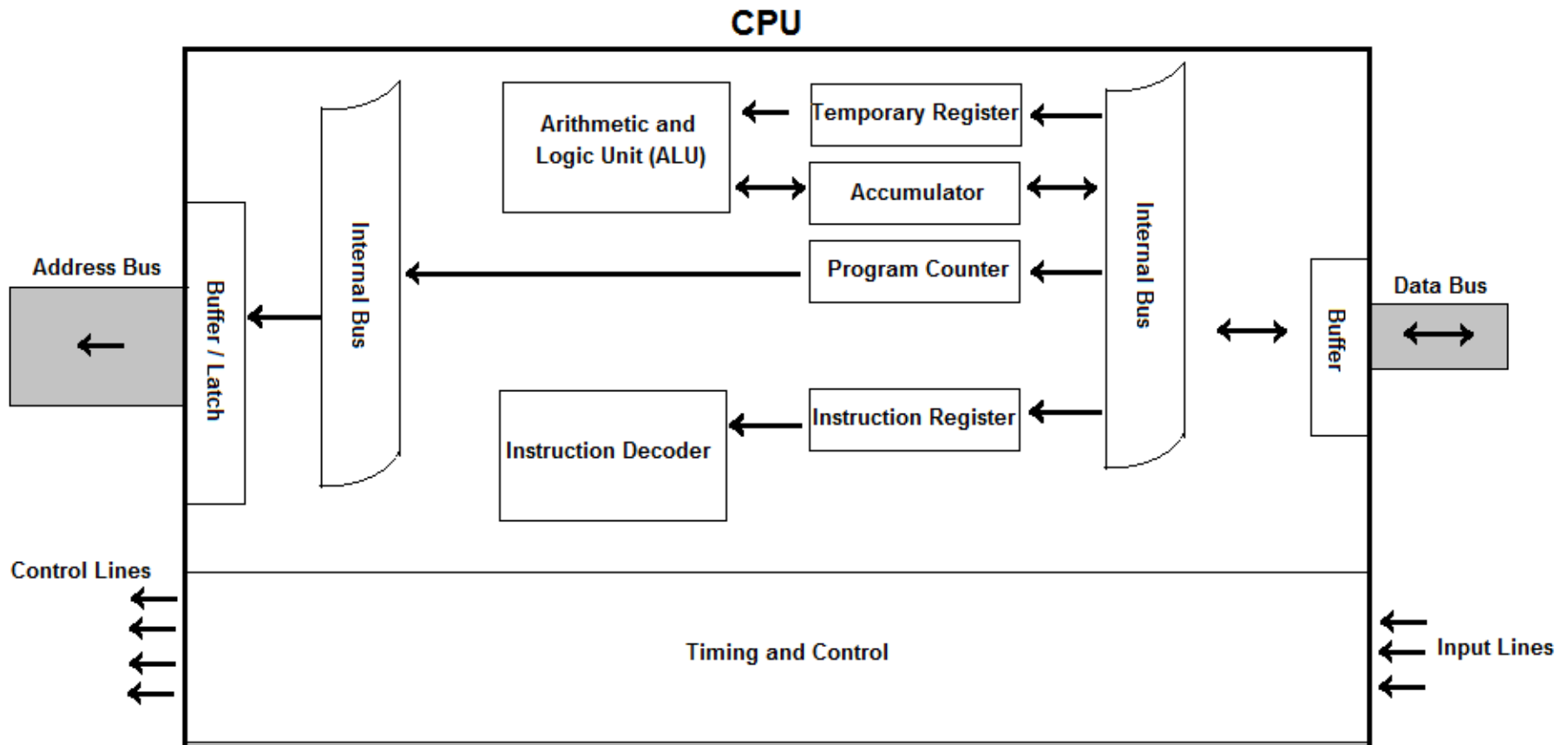
# Simplified CPU Architecture

- Microprocessor is the CPU
- CPU contains
  - Registers, computational circuitry, control circuitry
- Primary functions of CPU include:
  - Fetch, Decode and Execute instructions in proper order
  - Transfer data to and from memory and input/output
  - Respond to the interrupts
  - Provide Timing and control signals to other devices
- Various Registers
- The ALU
- Instruction Decoder
- Control and Timing Section
- Fetch Decode and Execute cycle



# Simplified CPU Architecture

- Architecture of generic microprocessor





# The Generic Microprocessor

- **Common characteristics of microprocessors**
  - Power connections
  - Bit size
  - Data lines
  - Address lines
  - Control lines
  - Internal registers
  - Addressing modes



# Data Sheet Description

- **Contains information like:**
  - IC packaging
  - Pin Diagram
  - Structure of microprocessor
  - Timing Diagram
  - Schematics of minimal system
  - Instruction set



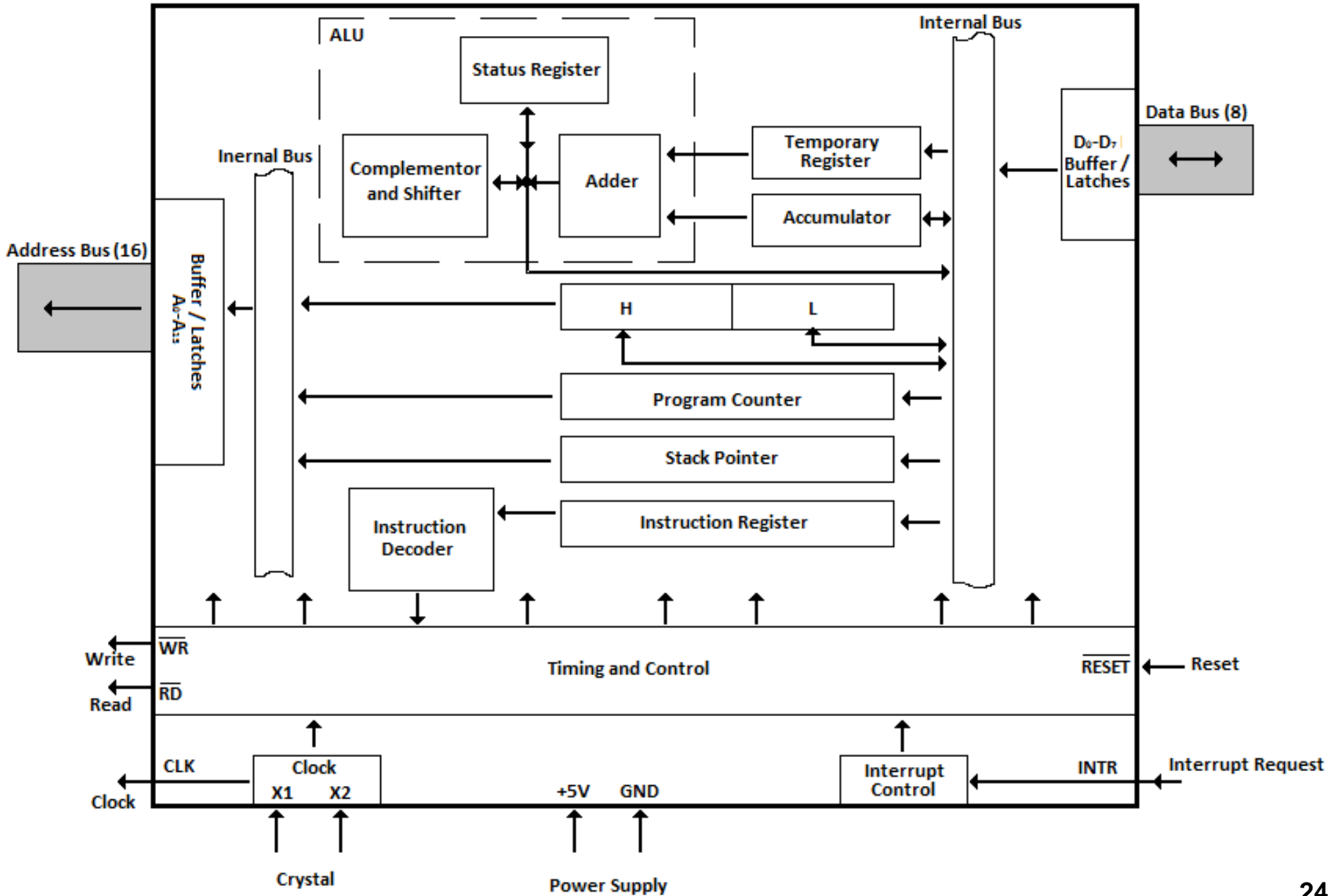
# Pin Diagram and Functions

- **40-pin DIP**
- **Regulated Power supply (1,2)**
- **Clock (38-40)**
- **16-bit address bus (5-20)**
- **8 bi-directional data lines (21-28)**
- **RD\*, WR\* control outputs (30,31)**
- **RESET (33)**
- **INTR (35)**
- **Steps in processing interrupts**



# CPU Architecture

- **Arithmetic and Logic Unit**
- **Several Registers**
- **Program Counter**
- **Instruction Decoder**
- **Bus Buffers and latches**
- **Control Lines**
- **T-States and Machine Cycles**
- **Microcontrollers**







# Using Data / Address Registers

- HL Register Pair
- It is a pointer to memory
- Simple program for adding three consecutive numbers in memory and storing in next location



# Using the Stack Pointer

- 16-bit special purpose counter-register that contains an address
- The address refer to the area of memory called stack
- Stack is used when there is a procedure call or interrupt
- Stack is also used to store temporarily the registers
- SP initialized at one location higher than the last word of stack
- Stack is read/write memory
- Data is written using PUSH or CALL instruction
- Data is read using POP or RETURN instruction
- It is a Last-In-First-Out Data Structure



# Using the Stack Pointer (Contd.)

- When reading the data from the stack, the data comes in reverse order
- When to POP data, first data is read and then SP is incremented
- PUSH and POP are always used in pair
- Every microprocessor has its own operand type to PUSH and POP



# Programming Generic Microprocessor

- A machine program is what a microprocessor can understand directly (composed of 1's and 0's)
- Program are difficult for human beings to understand
- The machine language program can be made easier to understand by replacing binary codes with Hexadecimal equivalents
- Still difficult to understand / program
- The program:
  - (1) Load the binary number (1011 0100) into accumulator
  - (2) Take 1's complement of the number
  - (3) Store the complemented number in memory location 2100h



# Programming Generic Microprocessor (Contd.)

- The last representation was easy and understandable by the humans
- One step up from the machine language is the “*assembly language*”
- Generally, the assembly language combines 1 or more bytes of machines code into a phrase or instruction
- Assembly language represents binary op-codes with *mnemonics*
- Like the mnemonic for the second instruction in former program is CMA (Complement Acc.) while its op-code is 0010 1111
- A program called *assembler* can translate mnemonics into binary codes (machine code)



# Programming Generic Microprocessor (Contd.)

- A program listing containing assembly language instructions, their machine codes and descriptions
- Such a listing contains several fields like address, machine code, label, mnemonic and operands
- The assembler only requires mnemonic and operands to translate
- The listing containing mnemonics and operands is called *source code* and is passed to assembler
- Assembler translates and produces an object code containing machine codes to be stored in consecutive memory locations



# Programming Generic Microprocessor

## (Contd.)

Address (hex)	Contents (hex)	Label	Op code	Operand	Comments
2000 2001 2002	21 20 20		LXI	H,2020H	Initialize HL register pair at 2020H
2003	06		MVI	B, 01H	Initialize register B with 01H
2004	01				
2005	7E		MOV	A, M	Move binary number from memory Location 2020H to accumulator
2006	2F		CMA		Complement accumulator
2007	80		ADD		Add B to A to form 2s complement
2008	77		MOV	M, A	Move 2s complement number from accumulator to memory location 2020H
2009	76		HLT		Stop MPU



# Questions

